This application note describes the implementation and use of a library that allows on-chip flash memory to be used for non-volatile storage, in a similar manner to EEPROM.
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1. Introduction

The NXP LPC1100, LPC1300, LPC1700 and LPC2000 family of ARM7 and ARM Cortex microcontrollers provide on-chip flash memory for storage of firmware. A frequent requirement for embedded systems is the storage of variables whose values need to persist through resets of the system. For microcontrollers without on-chip EEPROM there are two basic options. One is to attach an external EEPROM device, perhaps using I²C or SPI. The second is to use the on-chip flash memory for variable storage.

Attaching an external device requires additional integration work as well as increasing hardware production costs and board space. Using flash memory is problematic as flash memory is not arranged in an optimal way for storage of small pieces of changing data. There are three key issues with using flash memory for storage of variables:

- Erasing memory is on a per-sector basis
- Flash memory lifespan
- Reprogramming locations

This application note describes a simple solution to the problem of using flash memory for variable storage, which takes into account erasing flash memory in a way that avoids data loss and minimizes the effect of flash memory lifespan.

2. Storing variables in flash memory

When a flash memory location is erased its value becomes 0xFF. Each bit can then be programmed to logic 0 but it cannot be returned to logic 1 without erasing the entire sector. Subject to some limitations this allows multiple writes to a single flash sector to add data to it, without needing to erase the entire sector each time.

Each location in flash has a finite lifespan, rated in terms of the number of programming and erase cycles. For the Microcontrollers covered by this application note it is a minimum of 10,000 programming and erase cycles.

To continually erase a sector and program the same location each time a variable changed would result in continual use of the same location until the point of failure, while leaving thousands of other locations in the same sector unused.

The solution presented in this application note is to make use of all the locations in a single sector until it is full, then reorganize the memory usage so the process can start again.

2.1 Basic mechanism

Initially flash sector A is empty (i.e., erased). When a variable needs to be stored a variable record is programmed into sector A. Later when the variable needs to be updated a new record is programmed into sector A and the old record is ignored.

Eventually sector A will become full, mostly with old records that are no longer used. On the next variable write when there isn’t space in sector A, a second sector, sector B is employed.

All the current records in sector A are copied to sector B and sector A is then erased.

When a variable needs to be updated, the update goes into sector B, which is the sector currently in use. Once sector B has been filled the current records are copied to sector A and the process repeats.
By swapping between two sectors of the same size it can be ensured that current variable values are always stored somewhere in flash memory while erasing takes place. This helps to ensure that if a reset or power failure occurs; the risk of data loss is minimized.

![Diagram of basic mechanism](image)

**Fig 1. Basic mechanism**

### 2.2 Variable records

The following table describes the structure of a variable record.

<table>
<thead>
<tr>
<th>Field</th>
<th>Size in Bytes</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Flags</td>
<td>1</td>
<td>0xFF = no entry, 0xAA = valid data</td>
</tr>
<tr>
<td>ID</td>
<td>2</td>
<td>A unique identifier for the variable</td>
</tr>
<tr>
<td>Data</td>
<td>N</td>
<td>The data stored in native-endianess, maximum of 12 bytes</td>
</tr>
<tr>
<td>Padding</td>
<td>12 – N</td>
<td>Always set to 0x00</td>
</tr>
<tr>
<td>Checksum</td>
<td>1</td>
<td>A 2's compliment checksum of the ID, data and padding</td>
</tr>
</tbody>
</table>

Each record is always a multiple of 16 bytes due to a system called Error Correction Code (ECC) which can be found on some NXP microcontrollers. ECC allows one write per 16 consecutive bytes of flash memory.

If needed the record size could be increased, which would allow larger variables to be used at the expense of the number of variables that can be stored, however it must always be a multiple of 16 bytes. It is also possible to store more complex data types, such as structures, in a variable record.
The flags field indicates the status of the record. Erased memory has the effect of setting the flags to 0xFF which indicates no entry. This field takes advantage of the fact that it is possible to continually reprogram a location in flash memory without erasing, providing that each successive write only changes bits from logic 1 to logic 0. Therefore the lifecycle of a record is: no entry to valid data.

Because multiple variables are supported, it must be possible to identify the variable that a record contains. Therefore a unique identifier is used. For example if a variable called Counter is used, the ID might be defined using the C preprocessor:

```
#define COUNTERID 1
```

The same identifier must be used for a specific variable throughout the operation of the firmware.

The checksum is calculated by adding together the ID, data and padding using an eight byte variable, then subtracting from 0x100.

### 2.3 Sector organization

The start of both sectors used contains a single 48 byte sector record:

<table>
<thead>
<tr>
<th>Field</th>
<th>Size in Bytes</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Flags 1</td>
<td>1</td>
<td>0xFF = sector empty, 0xAA = sector not empty</td>
</tr>
<tr>
<td>Padding</td>
<td>15</td>
<td>Reserved. Always set to 0x00</td>
</tr>
<tr>
<td>Flags 2</td>
<td>1</td>
<td>0xFF = sector empty or initializing, 0xAA = sector valid or invalid</td>
</tr>
<tr>
<td>Padding</td>
<td>15</td>
<td>Reserved. Always set to 0x00</td>
</tr>
<tr>
<td>Flags 3</td>
<td>1</td>
<td>0xFF = sector empty, initializing or valid, 0xAA = sector invalid</td>
</tr>
<tr>
<td>Padding</td>
<td>15</td>
<td>Reserved. Always set to 0x00</td>
</tr>
</tbody>
</table>

The three flags fields indicate the status of the sector. An erased sector will have the effect of setting all the flags to 0xFF, which indicates the sector is empty. These fields take advantage of the fact that it is possible to continually reprogram a location in flash memory without erasing, providing that each successive write only changes bits from logic 1 to logic 0. The lifecycle of a sector is: empty to initializing to valid to invalid.

The flags are spaced 16 bytes apart. This is due to the ECC functionality previously mentioned, that is featured on some microcontrollers. By spacing the flags apart, any update to the flags results in a single write per 16 consecutive bytes.

The initializing state for a sector indicates that the sector is in the process of being filled with current variable records from the old sector. Once the copying has finished the sector will be marked as valid. When the sector has been filled and the current variable records have been copied to the new sector, it will be marked as invalid.

In the following diagram the flags for a sector have been grouped into 24-bit values with Flags 1 occupying the most significant byte, and Flags 3 the least significant byte.
2.4 Data access

To find the current value of a variable it is necessary to search through the current flash sector to find the last written valid record for the variable. If the value of the variable has changed then there will be several old records for the variable, which will need to be skipped. In the worst case the sector will be almost full and the variable record will be located at the end of the sector.

To speed up variable access a RAM-based lookup table is used. Every time a variable record is written an offset into the flash sector is stored in a table along with the variable’s ID. To read the value of a variable the lookup table is consulted and then a direct access to the variable record is made.

Similarly the next free location in the current sector is also stored in RAM, increasing the speed at which new variable records can be written.

Each entry in the RAM-based lookup table occupies eight bytes. As supplied the library supports a maximum of 100 variables (see #define MAX_VARIABLES), resulting in a
lookup table size of 800 bytes. If RAM usage is critical the RAM-based lookup table could be easily removed at the expense of slower non-volatile variable access.

2.5 **Recovery after reset**

The most important aspect of storing variables in flash memory is the ability to retain values even if the microcontroller is reset or power is interrupted. Even if power is interrupted while variables are being copied from one sector to another the system must be able to recover without data loss.

There is protection for this on two levels – variables level and sector level.

2.5.1 **Variable protection**

When the value for a new variable is stored in flash, the variable record has a flags field and a checksum field that ensure only complete and correctly written variable records are used. When reading the current value of a variable the flags must be set to 0xAA and the checksum must be correct.

After writing the variable record flash memory is read back to compare the contents and make sure the write was successful. If the write failed for some reason then the application layer is informed.

If the microcontroller is reset during the write of a variable record then after the next reset when the RAM-based lookup table is reconstructed, partially written or invalid records will be ignored. This will have the effect of using the last successfully written value for the variable.

2.5.2 **Sector protection**

During the switch between sectors the microcontroller could be interrupted at any point. After the following reset the microcontroller needs to continue where it left off switching sectors. To achieve this, a decision table is used. The flags of both sectors are examined to determine the current state of the system and what to do next. This takes place automatically during initialization. Note that not all states are possible during normal operation, however there has to be an action for every combination of sector flags.

In Table 3, the flags for a sector have been grouped into 24-bit values with Flags 1 occupying the most significant byte, and Flags 3 the least significant byte.

### Table 3. Sector decision table

<table>
<thead>
<tr>
<th>Sector A Flags</th>
<th>Sector B Flags</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>Invalid</td>
<td>X (don’t care)</td>
<td>Erase sector A</td>
</tr>
<tr>
<td>X (don’t care)</td>
<td>Invalid</td>
<td>Erase sector B</td>
</tr>
<tr>
<td>0xFFFFFF</td>
<td>0xFFFFFF</td>
<td>Mark sector A as valid (flags = 0xAAAAFF)</td>
</tr>
<tr>
<td>0xFFFFFF</td>
<td>0xAAAAFF</td>
<td>Mark sector B as valid (flags = 0xAAAAFF)</td>
</tr>
<tr>
<td>0xFFFFFF</td>
<td>0xAAAAAFF</td>
<td>No action</td>
</tr>
<tr>
<td>0xFFFFFF</td>
<td>0xAAAAAAA</td>
<td>Treat sector B as full and swap to sector A</td>
</tr>
<tr>
<td>0xAAFFFE</td>
<td>0xFFFFFF</td>
<td>Mark sector A as valid (flags = 0xAAAAFF)</td>
</tr>
<tr>
<td>0xAAFFFE</td>
<td>0xAAAAFF</td>
<td>Erase sector B. Mark sector A as valid (flags = 0xAAAAFF)</td>
</tr>
</tbody>
</table>
### Sector A Flags | Sector B Flags | Action
---|---|---
0xAAFFFF | 0xAAAAFF | Erase sector A. Treat sector B as full and swap to sector A
0xAAFFFF | 0xAAAAAA | Mark sector A as valid (flags = 0xAAAAFF). Erase sector B
0xAAAAFF | 0xFFFFFFF | No action
0xAAAAFF | 0xAAFFFF | Erase sector B. Treat sector A as full and swap to sector B
0xAAAAFF | 0xAAAAFF | Erase sector B
0xAAAAFF | 0xAAAAAA | Erase sector B
0xAAAAAA | 0xFFFFFF | Treat sector A as full and swap to sector B
0xAAAAAA | 0xAAAAFF | Mark sector B as valid (flags = 0xAAAAFF). Erase sector A
0xAAAAAA | 0xAAAAFF | Erase sector A
0xAAAAAA | 0xAAAAAA | Erase sectors A and B. Mark sector A as valid (flags = 0xAAAAFF)

### 2.5.3 Code Read Protection

All ARM-based microcontrollers from NXP feature Code Read Protection (CRP), which is a method of securing the device by programming special values into a predefined memory location.

The flash-based system described in this application note can function when the CRP is disabled or when it is set to CRP level one. However it will not function if the CRP is set to level two or three.

### 3. Application Programmer Interface

This section describes the interface to the library implementing the system of storing variables in flash memory.

#### 3.1 Functions

**3.1.1 NVOL_Init**

Initializes the non-volatile library. This function must be called before any other functions. Returns TRUE for success and FALSE for an error. An error can occur if the flash memory cannot be placed into a known state.

**3.1.2 NVOL_SetVariable**

Stores a value for a variable in flash memory. Passed is an ID unique to that variable along with a pointer to the variable and the size of the variable in bytes. This generic prototype allows for variables of any size and type to be stored. Returns TRUE for success and FALSE for an error.

**3.1.3 NVOL_GetVariable**

Gets the value of a variable previously stored in flash memory. Passed is an ID unique to that variable along with a location for the value to be stored and the size of the variable. This generic prototype allows for variables of any size and type to be read. Returns TRUE for success and FALSE for an error.
3.2 Configuration

Configuration options are set using the #define statements in flash_nvol.c.

3.2.1 Sector descriptions

The sector start addresses, numbers and sizes are set using SECTOR1_STARTADDR, SECTOR2_STARTADDR, SECTOR1_NUM, SECTOR2_NUM and SECTOR_SIZE. The sector number must match the numbering scheme given in the flash memory section of the NXP User Manual for the device.

3.2.2 CPU clock

The frequency of the CPU clock in kHz is set using CPU_CLK. This is needed by the IAP routines for timing.

3.2.3 Interrupts

Different compilers use different non-standard ways of enabling and disabling interrupts. The macros ENABLEIRQ and DISABLEIRQ must be set to the appropriate statements or function calls for the compiler you are using.

3.2.4 IAP access point

The value of IAP_LOCATION must be set to the memory location for IAP routines. This address can be found in the NXP User Manual of the device.

3.2.5 Number of variables and size

The maximum size of a variable in bytes is set using MAX_VARIABLE_SIZE. This must always be a multiple of 16 minus four (e.g., 12, 28, 44). A larger size results in less efficient usage of flash memory.

The maximum number of variables is set using MAX_VARIABLES. This has to be small enough to ensure that all variables can fit into a single flash sector. Note that 48 bytes are always allocated at the start of each sector for sector management. Also four bytes are needed per variable record. The smaller this number is, the less RAM is needed for the lookup table.

4. Limitations

There are two key limitations with the system. Firstly the flash lifespan limits the number of times that a variable can be written. Secondly there will be a delay in writing variables to flash memory.

4.1 Flash lifespan

The flash memory in the microcontrollers covered by this application note is rated at a minimum of 10,000 program and erase cycles.

If there is a single 12-byte variable and the flash sectors are 4096 bytes in size, then the variable can be written \((4096 - 16) / 12 = 340\) times before the sector is full. At that point the variable will move to the other sector where it can be written an additional 340 times. The result is that the variable can be written 680 times for a single erase in each sector. The lifespan is therefore a minimum of \(10,000 \times 680 = 6.8\) million variable writes.

Now consider the worst case for sectors 4096 bytes in size. It would be 340 variables of 12-bytes in size. Every write of a value would cause the flash sectors to be swapped,
resulting in one erase cycle of a specific sector every two variable writes. The lifespan would therefore be a minimum of 20,000 variable writes.

Note that the library only writes a variable if the value has changed. Consecutive writes with the same value will have no effect.

4.2 Write delays

Each time a new value for a variable is written it has to be programmed into flash memory. The speed of this operation is limited by the flash programming and erase speed.

The worst case write time for a single variable when the current sector is not full is approximately 1.05ms.

If the sector is full then the write will cause a swap of the two sectors used in the system. The worst case time for the entire operation will be approximately 105ms + 1.05ms per variable + 1.05ms for the current variable. For example if there are 20 variables then the approximate worst case time will be 105 + (1.05 x 20) + 1.05 = 127.05ms.

A possible optimization that could be implemented is to add a function that allows the application layer to determine how close the current sector is to being full and then manually perform a swap of the sectors. This would avoid unexpectedly long write delays at undesirable moments by ensuring that the sector swapping always takes place during idle periods.

4.3 IAP RAM usage

Calls to erase and program the flash memory using IAP require RAM. The location and amount of RAM used is defined in the NXP User Manual for the device, but is typically 32 bytes at the top of the RAM space and a maximum of 128 bytes on the user stack.
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